

LEZIONE 30/03

RISOLUZIONE DI SISTEMI LINEARI

PROBLEMA COMPUTAZIONALE:

Dato $A \in \mathbb{R}^{m \times m}$, $b \in \mathbb{R}^m$ determinare $x \in \mathbb{R}^m$:

$$\boxed{Ax = b}$$

A matrice dei coefficienti, b termine noto
 x vettore delle incognite.

Assunzione sui dati: $\det(A) \neq 0$

$$\Rightarrow \exists! x: Ax = b \quad \boxed{x = A^{-1}b}$$

① A è una matrice diagonale

$$A = (a_{ij}) \quad a_{ij} = 0 \text{ se } i \neq j$$

$$A = \begin{bmatrix} a_1 & & 0 \\ & \ddots & \\ 0 & & a_m \end{bmatrix}$$

$$\det A = \prod_{i=1}^m a_i$$

A invertible $\Leftrightarrow a_{ii} \neq 0 \quad i=1 \dots n$

$Ax = b \Leftrightarrow \begin{cases} a_{11}x_1 = b_1 \Leftrightarrow x_1 = b_1/a_{11} \\ \vdots \\ a_{nn}x_n = b_n \Leftrightarrow x_n = b_n/a_{nn} \end{cases}$

for $k=1:n$

$x(k) = b(k) / a(k,k);$

end.

↳ it's computer

$O(n)$ operation.

for $k=n:-1:1$

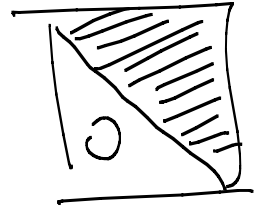
$x(k) = b(k) / a(k,k);$

end

• A is triangular...

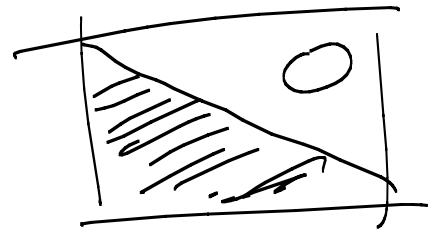
$A = (a_{ij})$ si dice TRIANGOLARE SUPERIORE

se $a_{ij} = 0$ per $i > j$



$A = (a_{ij})$ si dice TRIANGOLARE INFERIORE

se $a_{ij} = 0$ per $j > i$



$$A = \begin{pmatrix} a_{11} & \dots & a_{1n} \\ \vdots & \ddots & \vdots \\ 0 & \dots & a_{nn} \end{pmatrix}$$

$$\det A = \prod_{i=1}^n a_{ii}$$

A invertibile $\Leftrightarrow a_{ii} \neq 0 \quad (i=1, \dots, n)$

$Ax = b \quad (\Rightarrow)$

$$\begin{cases} a_{11}x_1 + \dots + a_{1n}x_n = b_1 \\ a_{21}x_1 + \dots + a_{2n}x_n = b_2 \\ a_{31}x_1 + \dots + a_{3n}x_n = b_3 \\ \vdots \\ a_{nn}x_n = b_n \end{cases}$$

BACKWARD SUBSTITUTION (SOSTITUZIONE ALL'INDIETRO)

$$a_{kk}x_k + a_{k,k+1}x_{k+1} + \dots + a_{kn}x_n = b_k$$

$$a_{kk}x_k + \sum_{j=k+1}^n a_{kj}x_j = b_k$$

$$x_k = \frac{1}{a_{kk}} \left(b_k - \sum_{j=k+1}^n a_{kj}x_j \right)$$

$k = n-1, \dots, 1$

$$x(n) = b(n)/a(n,n);$$

for $k = n-1, \dots, 1$

$$s = 0;$$

for $j = k+1, \dots, n$

$$s = s + a(k,j) * x(j);$$

end;

$$x(k) = (b(k) - s) / a(k, k);$$

bed

Costs computat.:

$$\sum_{k=1}^{n-1} (n-k) = \frac{(n-1) + (n-2) + \dots + 1}{2}$$

$$= \frac{n \cdot (n-1)}{2} \quad (\text{arithmetische \& INBOLORE})$$

$$= \frac{n^2}{2} + O(n) \quad \text{op.: multiplizieren}$$

$$= \underline{O(n^2)} \quad \text{op.: multiplizieren}$$

$$A \prec \begin{pmatrix} \diagdown & & \\ & \diagdown & \\ & & \diagdown \end{pmatrix} = \begin{pmatrix} a_{11} & & & \\ & \ddots & & \\ & & \ddots & 0 \\ & & & \ddots & \\ a_{n1} & & & & a_{nn} \end{pmatrix}$$

$$\Rightarrow \begin{cases} a_{11}x_1 = b_1 \\ \vdots \\ a_{n1}x_1 + \dots + a_{nn}x_n = b_n \end{cases}$$

FORWARD SUBSTITUTION (SOSTITUZIONE IN AVANTI)

~~→~~ IMPLEMENTAZIONE (ANALOGA)

$$A \in \mathbb{R}^{m \times n} \quad \text{genera} \quad A = \begin{bmatrix} a_{11} & \dots & a_{1n} \\ \vdots & & \vdots \\ a_{m1} & \dots & a_{mn} \end{bmatrix}$$
$$Ax = b$$

TRASFORMARE IL SISTEMA LINEARE IN UN SISTEMA EQUIVALENTE CON MATRICE DEI COEFFICIENTI PIÙ SEMPLICE (AD ESEMPIO ~~DIAGONALE~~ O TRIANGOLARE)

- Ognora A si può fattorizzare come prodotto di matrici TRIANGOLARI

$$A = L \cdot U \quad \text{A invertibile} \Rightarrow L, U \text{ invertibili}$$

$$Ax = b \Leftrightarrow \begin{matrix} L \\ U \end{matrix} x = b \Leftrightarrow \begin{cases} Ly = b \\ Ux = y \end{cases}$$

Def: L is a left ideal: Rows of $A \in \mathbb{F}^{m \times n}$

L is a subspace of \mathbb{F}^m and $\neq \{0\}$ unless trivial
 In fact we can think of L as a subspace of \mathbb{F}^m spanned by the

$$A \in \mathbb{F}^{m \times n}$$

$$A = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} u_1 & u_2 \\ 0 & u_3 \end{pmatrix}$$

$$\left. \begin{array}{l} u_1 = 0 \\ u_2 = 1 \\ u_1 = 1 \\ u_2 + u_3 = 0 \end{array} \right\} \Rightarrow \text{rows are } \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

$$A = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} u_1 & u_2 \\ 0 & u_3 \end{pmatrix}$$

$$\left. \begin{array}{l} u_1 = 1 \\ u_2 = 1 \\ u_1 = 1 \\ u_2 + u_3 = 1 \end{array} \right\} \Rightarrow \begin{array}{l} L = \mathbb{F} \\ u_3 = 0 \end{array}$$

∃ infinite families LU

Thm: So $A \in K^{n \times n}$ is row

$$A_{kk} = A_{(1:k, 1:k)}, \quad 1 \leq k \leq n,$$

be the submatrix principle of A of order $1, 2, \dots, n$.

$$\text{If } \det(A_k) \neq 0 \quad \underline{\underline{k=1 \dots n-1}}$$

then $\exists!$ LU decomposition of A .

Pr: Induction on n (dimension of A)

$$n=1 \quad \underline{A} = [a]$$

$$a = 1 \cdot x \Leftrightarrow x = a$$

$$\underline{a = 1 \cdot a}$$

Ipten' uolutas: Assumon el teom cas fno el'atun m
 Le $B \in K^{(n-1) \times (n-1)}$ e B_1, \dots, B_{n-2} sse matrob:
 oler \underline{I} ! f'ltan' LU el B

$u.$ $A \in \mathbb{R}^{n \times n}$ $A = \left[\begin{array}{c|c} A_{n-1} & z \\ \hline v^T & a \end{array} \right]$

$$\left[\begin{array}{c|c} A_{n-1} & z \\ \hline v^T & a \end{array} \right] = \left[\begin{array}{c|c} L_{n-1} & 0 \\ \hline x^T & 1 \end{array} \right] \left[\begin{array}{c|c} U_{n-1} & y \\ \hline 0 & \beta \end{array} \right]$$

$$\Leftrightarrow \begin{cases} L_{n-1} U_{n-1} = A_{n-1} & (\text{f'ltan' LU el } A_{n-1}) \\ L_{n-1} y = z & \Leftrightarrow y = L_{n-1}^{-1} z \\ x^T U_{n-1} = v^T & \Leftrightarrow x^T = v^T U_{n-1}^{-1} \\ x^T y + \beta = a & \Leftrightarrow \beta = a - x^T y \end{cases}$$

$A_{n-1} \in K^{(n-1) \times (n-1)}$ per' ipten' el' teom. A_1, \dots, A_{n-2}
 sse matrob. \Rightarrow per' ipten' uolutas $\Rightarrow \underline{I}$!

Induziert LU of A_{m-1} $A_{m-1} = \underline{\underline{L_{m-1}}} \underline{\underline{U_{m-1}}}$

$$\det(A_{m-1}) = \det U_{m-1} \neq 0$$

\Rightarrow $\underline{\underline{A}}$ ed \bar{c} univariate system \Rightarrow $\underline{\underline{A}}$ LU \Rightarrow $\underline{\underline{A}}$

$$A_2 \left(\begin{array}{ccc|c} 1 & 2 & 3 & \\ \hline 4 & 5 & 6 & \\ \hline 7 & 8 & 9 & \end{array} \right)$$

$$A_1 = [1]$$

$$A_3 = A$$

$$A_2 = \begin{bmatrix} 1 & 2 \\ 4 & 5 \end{bmatrix}$$

$$x^T = [x_1 \dots x_{m-1}] \quad x = \begin{bmatrix} x_1 \\ \vdots \\ x_m \end{bmatrix}$$

$$A_2 \left(\begin{array}{ccc|c} 1 & & & x_1 \\ & \ddots & & \vdots \\ & & 1 & x_m \\ \hline x_1 & \dots & & \end{array} \right)$$

$$\in \mathbb{R}^{m \times m}$$

(freew.)
arrow wheel
arrow wheel



- ① Dăruiește o anumită factorizare LU
- ② În caz afirmativ, determină toate factorizările.
- ③ Determină, pentru o dată x la un a și ϵ singur:
- ④ Scrie un program Matlab pentru calculul lui $Ax=b$ în cazul în care L are singurii nenuli explicit, toate resturile.

① $A \in \mathbb{R}^{n \times n}$ $A_k = \begin{pmatrix} 1 & 0 \\ 0 & \alpha \end{pmatrix} \dots A_k = I$ și $\alpha \neq 0$
 $k=1, \dots, n-1$

$\Rightarrow \det A_k = \alpha \neq 0 \quad k=1, \dots, n-1$

$\Rightarrow \exists$! factorizare LU de A .

②
$$\left[\begin{array}{c|c} I & \begin{matrix} x_1 \\ \vdots \\ x_{n-1} \end{matrix} \\ \hline x_1, \dots, x_{n-1} & x_n \end{array} \right] = \left[\begin{array}{c|c} I_{n-1} & 0 \\ \hline v^T & \alpha \end{array} \right] \left[\begin{array}{c|c} I_{n-1} & z \\ \hline 0 & \beta \end{array} \right]$$

$L_{n-1} U_{n-1} = I \Rightarrow L_{n-1} = I \quad U_{n-1} = I$

$$Z = \begin{bmatrix} x_1 \\ \vdots \\ x_{n-1} \end{bmatrix} \quad v^T = [x_1 \dots x_{n-1}]$$

$$v^T Z + \beta = x_n \Leftrightarrow \beta = x_n - (x_1^2 + x_2^2 + \dots + x_{n-1}^2)$$

$$Az = \left(\begin{array}{c|c} I_{n-1} & 0 \\ \hline x_1 \dots x_{n-1} & 1 \end{array} \right) \left(\begin{array}{c|c} I_{n-1} & x_1 \\ \hline 0 & \beta \end{array} \right)$$

$$\textcircled{3} \text{ set } A = 0 \Leftrightarrow \text{set } U = 0 \Leftrightarrow$$

$$\beta = 0 \Leftrightarrow x_n = x_1^2 + x_2^2 + \dots + x_{n-1}^2$$

$$\textcircled{4} Az = b \Leftrightarrow LUz = b \Leftrightarrow \begin{cases} Ly = b \\ Uz = y \end{cases}$$

$$\equiv \begin{bmatrix} 1 & & & & & \\ & \ddots & & & & \\ & & \ddots & & & \\ & & & \ddots & & \\ & & & & \ddots & \\ & & & & & 1 \end{bmatrix} \begin{bmatrix} y_1 \\ \vdots \\ y_n \end{bmatrix} = \begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix} \quad \equiv \text{so that we can write}$$

$$\begin{array}{c} ||| \\ \left[\begin{array}{c|c} 1 & \\ \hline & 0 \end{array} \right] \begin{array}{c} x_1 \\ \vdots \\ x_{n-1} \\ 0 \end{array} \left[\begin{array}{c} z_1 \\ \vdots \\ z_m \end{array} \right] \left[\begin{array}{c} y_1 \\ \vdots \\ y_m \end{array} \right]
 \end{array}$$

Back substitution

$$\left[\begin{array}{c|c} 1 & \\ \hline x_1 \dots x_{m-1} & 1 \end{array} \right] \begin{array}{c} y_1 \\ \vdots \\ y_m \end{array} = \begin{array}{c} b_1 \\ \vdots \\ b_m \end{array} \quad (=)$$

$$\left\{ \begin{array}{l} y_1 = b_1 \\ y_2 = b_2 \\ \vdots \\ y_{m-1} = b_{m-1} \\ x_1 y_1 + \dots + x_{m-1} y_{m-1} + y_m = b_m \end{array} \right.$$

$$y(n:m-1) = b(n:m-1); \\
 s = 0;$$

for $k = n-1$

$$s = s + x(k) * y(k);$$

$O(m)$ operations

$$\text{end} \\
 y(m) = b(m) - s$$

$$\begin{bmatrix} 1 & & & & & \\ & \ddots & & & & \\ & & \ddots & & & \\ & & & \ddots & & \\ & & & & \ddots & \\ & & & & & \beta \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ x_{m-1} \\ x_m \end{bmatrix} = \begin{bmatrix} z_1 \\ \vdots \\ z_{m-1} \\ z_m \end{bmatrix} = \begin{bmatrix} y_1 \\ \vdots \\ y_m \end{bmatrix}$$

$\beta = x_m - (x_1 + \dots + x_{m-1})$

$$\beta = x_m$$

for $k = 1 : m-1$

$$\beta = \beta - x(k) * x(k);$$

$(0 \text{ } \infty)$

$$z(m) = y(m) / \beta;$$

for $k = m-1 : -1 : 2$

$$z(k) = y(k) - x(k) * z(m);$$

$(0 \text{ } \infty)$

end

$$A = \begin{bmatrix} 1 & & & & \\ & \ddots & & & \\ & & \ddots & & \\ & & & \ddots & \\ & & & & \ddots \\ & & & & & \beta \end{bmatrix} \in \mathbb{R}^{m \times m}$$

① Zeigen & Annahme für LU existenz von offener Ordnung
 der Matrizen

② Der prinzipielle Algorithmus für LU

$$A \in \mathbb{R}^{n \times n} \quad k=1, \dots, n-1 \quad \Rightarrow \quad A_k \text{ invertierbar}$$

$$\Rightarrow \exists! \text{ für Matrizen } \underline{LU}$$

$$A_k = \left(\begin{array}{ccc|c} \uparrow & & & \alpha \\ \uparrow & \downarrow & & \\ \uparrow & & \downarrow & \\ \hline & & & \beta \end{array} \right) = \left(\begin{array}{ccc|c} \uparrow & & & 0 \\ \uparrow & \downarrow & & \\ \uparrow & & \downarrow & \\ \hline & & & \beta \end{array} \right)$$

$$L_{n-1} U_{n-1} = \left(\begin{array}{ccc} \uparrow & & \\ \uparrow & \downarrow & \\ \uparrow & & \downarrow \end{array} \right) = \left(\begin{array}{ccc} \uparrow & & \\ \uparrow & \downarrow & \\ \uparrow & & \downarrow \end{array} \right) I$$

$$\left(\begin{array}{ccc} \uparrow & & \\ \uparrow & \downarrow & \\ \uparrow & & \downarrow \end{array} \right) \begin{pmatrix} z_1 \\ \vdots \\ z_m \end{pmatrix} = \alpha \cdot \begin{pmatrix} \beta \\ \vdots \\ 0 \end{pmatrix}$$

$$\begin{cases} z_1 = \alpha \\ z_1 + z_2 = 0 \\ z_1 + z_2 + z_3 = 0 \\ \vdots \\ z_1 + z_2 + \dots + z_n = 0 \end{cases}$$

$$Z_{12} = 1 \quad Z_{22} = -1 \quad Z_3 = Z_4 = \dots = Z_n = 0$$

$$\cancel{\alpha} - \cancel{\alpha} + \beta = 1$$

$$\Leftrightarrow \beta = 1$$

$$A = \left(\begin{array}{c|c} \begin{array}{c} \uparrow \\ \downarrow \end{array} & \begin{array}{c} \uparrow \\ \downarrow \end{array} \\ \hline \begin{array}{c} \uparrow \\ \downarrow \end{array} & \begin{array}{c} \uparrow \\ \downarrow \end{array} \end{array} \right) = \left(\begin{array}{c|c} \begin{array}{c} \uparrow \\ \downarrow \end{array} & \begin{array}{c} \uparrow \\ \downarrow \\ \alpha \\ \alpha \\ 0 \\ \dots \\ 0 \\ \alpha \end{array} \\ \hline \begin{array}{c} \uparrow \\ \downarrow \end{array} & \begin{array}{c} \uparrow \\ \downarrow \end{array} \end{array} \right) = \left(\begin{array}{c|c} \begin{array}{c} \uparrow \\ \downarrow \end{array} & \begin{array}{c} \uparrow \\ \downarrow \end{array} \\ \hline \begin{array}{c} \uparrow \\ \downarrow \end{array} & \begin{array}{c} \uparrow \\ \downarrow \end{array} \end{array} \right)$$

$$\det A = \det(L \cdot U) = \det L \cdot \det U$$

$$= \det U = 1$$

$\forall \alpha \in \mathbb{R}$ la matrice est inversible.

$$\det \begin{pmatrix} \alpha & 0 & 0 & \alpha \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 \end{pmatrix} = 1 - \alpha \det \begin{pmatrix} 1 & 1 & 0 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix} = 1$$

Def: $A \in K^{m \times n}$ is dice predominant through 'per' ugle.

$$\text{se } |a_{ii}| > \sum_{\substack{j=1 \\ j \neq i}}^m |a_{ij}| \quad i=1, \dots, m$$

Def: $A \in K^{m \times n}$ is dice predominant through 'per' abou.

$$\text{se } |a_{ii}| > \sum_{\substack{j=1 \\ j \neq i}}^m |a_{ji}| \quad i=1, \dots, m$$

$$A = \begin{pmatrix} 3 & -1 & 1 \\ -1 & 5 & -3 \\ 1 & 1 & 3 \end{pmatrix}$$

$$|3| > 2 = |1| + |-1|$$

$$|5| > |-1| + |-3| = 4$$

$$|3| > |-1| + |1| = 2$$

Thm: So $A \in K^{m \times n}$ predominant through 'per' ugle.

Alor A is invertible.

Thm: Suppose A is base of rank n .

So $A \in K^{n \times n}$ predominant through 'per' ugle \Rightarrow

$$0 \notin K_i \quad (i=1, \dots, m) \Rightarrow 0 \notin \bigcup K_i$$

\Rightarrow 0 non è autovalore di $A \Leftrightarrow A$ è invertibile.

$$K_i = \left\{ z \in \mathbb{C} : |z - a_{ii}| \leq \sum_{\substack{j=1 \\ j \neq i}}^m |a_{ij}| \right\}$$

$$\underline{0 \notin K_i?} \quad |0 - a_{ii}| = |a_{ii}| \stackrel{\text{p.d.}}{>} \sum_{\substack{j=1 \\ j \neq i}}^m |a_{ij}|$$

$$\Rightarrow 0 \notin K_i \quad \square$$

Se A p.d. p.d. b. autovalori, certi p.d.:

$$\begin{bmatrix} 0 & 2 \\ 1 & 0 \end{bmatrix}$$

$$\left[\begin{array}{l} A \text{ p.d.} \Rightarrow A_x \text{ è p.d.} \Rightarrow \\ \exists ! \text{ forma LU di } A. \end{array} \right.$$

$$A = \begin{bmatrix} 3 & & & & & \\ & 1 & & & & \\ & & \ddots & & & \\ & & & 0 & & \\ & & & & \ddots & \\ & & & & & 1 \\ & & & & & & 2 \\ & & & & & & & 3 \end{bmatrix}$$

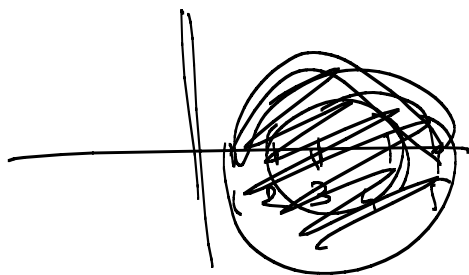
① A ammette f.f. LU?

① A 'p.d.' $3 = |3| > |1| = 2 \quad |z_{1,m}|$

$3 = |3| > |1| + |1| = 2 \quad 2 \leq i \leq n-1$

$\Rightarrow \exists!$ f.f. LU di A

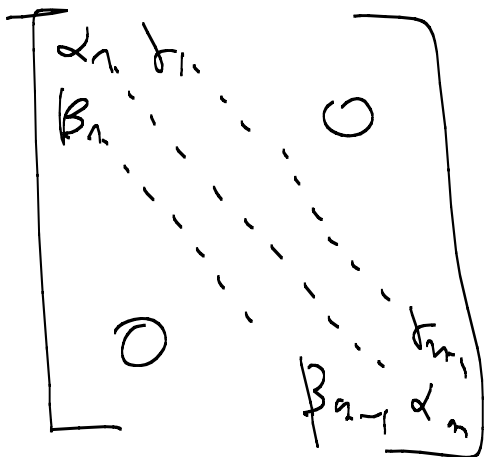
② Jacobian $A_{k \times k} = \begin{pmatrix} 3 & & & \\ & 1 & & \\ & & \ddots & \\ & & & 0 \\ & & & & \ddots \\ & & & & & 1 \\ & & & & & & 2 \\ & & & & & & & 3 \end{pmatrix}$



$\exists \frac{m}{k} \Rightarrow A_{k \times k}$ invertibile

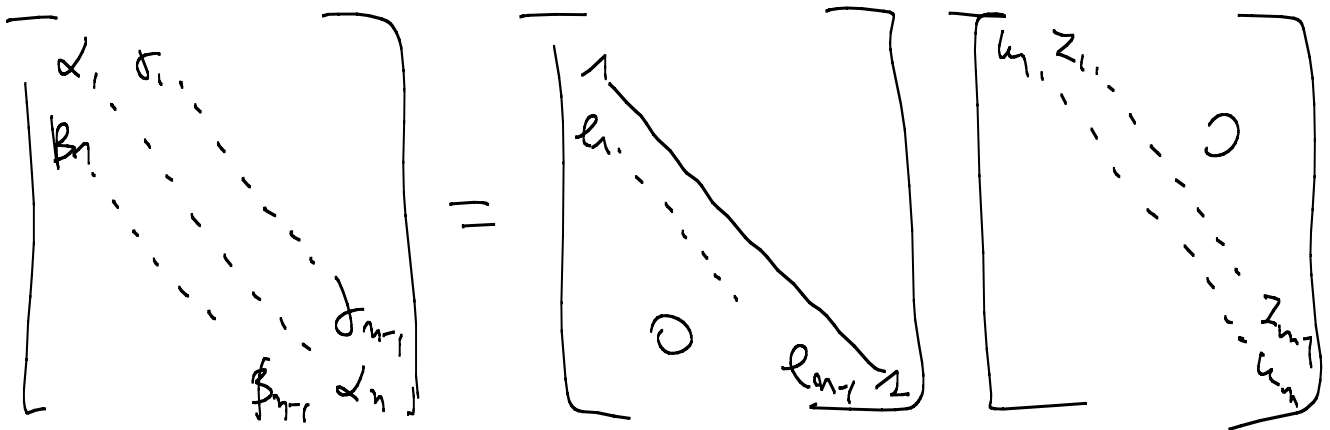
\Rightarrow Teorema di esistenza e unicità $\exists!$ f.f. LU

Calcolo veloce FATTORIZZAZIONE TRU COMPICATA
(ALGORITMO THOMAS)



matrice tridiagonală.

Suprasau de cantitate de LU



$$u_{12} \alpha_2 \quad Z_{12} \gamma_3$$

$$l_{12} u_{12} \beta_1 \quad (z) \quad l_{12} = \beta_{12} / u_{12}$$

$$l_{12} z_1 + u_{22} \alpha_2 \quad \Leftrightarrow \quad u_{22} \alpha_2 - l_{12} z_1 \quad Z_{12} \gamma_2$$

$$u(k) \alpha(k);$$

for $k = 2 : m$

$$Z(k-1) \alpha(k);$$

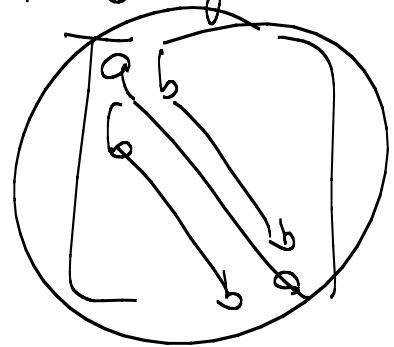
$$l(k) = \beta(k-1) / u(k-1)$$

$$u(k) = d(k) - l(k-1) * z(k-1);$$

but

$O(m)$ operations

Exercise: Implement in MATLAB the algorithm of Thoms
applied to a matrix in the step



INPUT: (a, b, m)

Output: l, u, z (column)

• Verify yourself the correctness of the algorithm of Thoms
